

The Gödel fibration

Davide Trotta

j.w.w. M. Spadetto and V. de Paiva

University of Pisa

5-2021



Introduction: Dialectica interpretation

Gödel's Dialectica Interpretation: an interpretation of intuitionistic arithmetic HA in a quantifier-free theory of functionals of finite type, called system \mathbb{T} .

Idea: translate every formula A of HA to $A^D = \exists x \forall y A_D$, where A_D is quantifier-free.

Application: if HA proves A , then system \mathbb{T} proves $A_D(t, y)$, where y is a string of variables for functionals of finite type, and t a suitable sequence of terms (not containing y).

Goal: to be as constructive as possible, while being able to interpret all of classical arithmetic.

Gödel (1958), *Über eine bisher noch nicht benützte erweiterung des finiten standpunktes*, Dialectica, 12(3-4):280–287.

Introduction: Dialectica interpretation

Dialectica category: given a category \mathcal{C} with finite limits, one can build a new category $\mathcal{D}\text{ial}(\mathcal{C})$, the objects of which have the form (X, U, α) where α is a subobject of $X \times U$ in \mathcal{C} ; such an object is thought of as the formula

$$\exists x \forall u \alpha(x, u).$$

An arrow from $\exists x \forall u \alpha(x, u)$ to $\exists y \forall v \beta(y, v)$ can be thought of as a pair (f, f_0) of terms, subject to the condition

$$\alpha(x, f_0(x, v)) \vdash \beta(f(x), v).$$

The definition of morphism is motivated by the way the dialectica interpretation acts on implicational formulae.

Introduction: Dialectica interpretation

Dialectica pseudo-monad: given a fibration p , one can construct the Dialectica fibration $\mathfrak{D}ial(p)$. Moreover, under the assumption that the base category of p is cartesian closed, this construction is monadic.

Hyland (2002), *Proof theory in the abstract*, *Annals of Pure and Applied Logic*, 114(1):43 - 78
Hofstra (2011), *The dialectica monad and its cousins*, *Models, logics, and higherdimensional categories: A tribute to the work of Mihály Makkai*, 53:107-139

Our contributions

- Given a fibration p , when is there a fibration p' such that $\mathcal{D}ial(p') \cong p$?
- When such fibration p' exists, how is it done?

Background

Definition

Let $p: E \longrightarrow B$ be a functor and $X \xrightarrow{f} Y$ an arrow in E . Let us call $A \xrightarrow{u:=p(f)} B$ the arrow $p(f)$ of B . We say that f is **Cartesian over u** if, for every morphism $Z \xrightarrow{g} Y$ in E such that $p(g)$ factors through u , $p(g) = uw$, there exists a unique $Z \xrightarrow{h} X$ of E such that $g = fh$ and $p(h) = w$.

Definition

A **fibration** is a functor $p: E \longrightarrow B$ such that, for every Y in E and every $I \xrightarrow{u} pY$, there exists a Cartesian arrow $X \xrightarrow{f} Y$ over u .

Jacobs (1999), *Categorical Logic and Type Theory*, Studies in Logic and the foundations of mathematics, 141

Background

Definition

We say a fibration $p: E \longrightarrow B$ over a category B with finite products has **simple coproducts** when the weakening functors π^* have left adjoints \coprod_{π} satisfying the *Beck-Chevalley Condition* (abbreviated as BCC).

Dually, we say that a fibration $p: E \longrightarrow B$ has **simple products** when the weakening functors π^* have right adjoints \prod_{π} satisfying BCC.

Quantifier-free objects

The logical intuition behind the next definition is that an element α is *quantifier-free* if it satisfies the following universal property: if there is a proof π of a statement $\exists i \beta(i)$ assuming α , then there exists a *witness* t , which depends on the proof π , together with a proof of $\beta(t)$. Moreover, we require that this holds for every re-indexing $\alpha(f)$ because in logic quantifier-free propositions are stable under substitution.

Quantifier-free objects

Definition

Let $p: E \rightarrow B$ be a fibration with simple coproducts. An object α of the fibre E_I is said to be **\coprod -quantifier-free** if it enjoys the following universal property: for every pair of arrows

$$A \times B \xrightarrow{\pi_A} A \xrightarrow{f} I$$

and every vertical arrow:

$$f^* \alpha \xrightarrow{h} \coprod_{\pi_A} \beta$$

of E_A , where β is an object of the fibre $E_{A \times B}$, there exist a unique arrow $A \xrightarrow{g} B$ of B and a unique vertical arrow $f^* \alpha \xrightarrow{\bar{h}} \langle 1_A, g \rangle^* \beta$ of E_A such that:

$$h = \left(f^* \alpha \xrightarrow{\bar{h}} \langle 1_A, g \rangle^* \beta \xrightarrow{\langle 1_A, g \rangle^* \eta_\beta} \langle 1_A, g \rangle^* \left(\pi_A^* \coprod_{\pi_A} \beta \right) = \coprod_{\pi_A} \beta \right)$$

Quantifier-free objects

Definition

We say that a fibration with simple coproducts $p: E \rightarrow B$ has **enough \coprod -quantifier-free objects** if, for every object I of B and for every element $\alpha \in E_I$, there exist an object A and a \coprod -quantifier-free object β in $E_{I \times A}$ such that $\alpha \cong \coprod_{\pi_I} \beta$.

Quantifier-free objects

Definition

Let $p: E \rightarrow B$ be a fibration with simple products. An object α of the fibre E_I is said to be **\prod -quantifier-free** if it enjoys the following universal property: for every arrow f and every projection π_A in B as follows:

$$A \times B \xrightarrow{\pi_A} A \xrightarrow{f} I$$

and every vertical arrow:

$$\prod_{\pi_A} \beta \xrightarrow{h} f^* \alpha$$

of E_A , where β is an object of the fibre $E_{A \times B}$, there exist a unique arrow $A \xrightarrow{g} B$ of B and a unique vertical arrow $\langle 1_A, g \rangle^* \beta \xrightarrow{\bar{h}} f^* \alpha$ of E_A such that:

$$h = \left(\prod_{\pi_A} \beta = \langle 1_A, g \rangle^* \left(\pi_A^* \prod_{\pi_A} \beta \right) \xrightarrow{\langle 1_A, g \rangle^* \varepsilon_\beta} \langle 1_A, g \rangle^* \beta \xrightarrow{\bar{h}} f^* \alpha \right)$$

Quantifier-free objects

Definition

We say that a fibration with simple products $p: E \rightarrow B$ has **enough- \prod -quantifier-free objects** if, for every object I of B and for every element $\alpha \in E_I$, there exist an object A and a \prod -quantifier-free object β in $E_{I \times A}$ such that $\alpha \cong \prod_{\pi_I}(\beta)$.

Skolem fibration

Definition

A fibration $p: E \rightarrow B$ is called a **Skolem fibration** if:

- its base category B is cartesian closed;
- the fibration p has simple products and simple coproducts;
- the fibration p has enough \prod -quantifier-free objects.
- \prod -quantifier-free objects are stable under simple products, i.e. if $\alpha \in E_I$ is a \prod -quantifier-free object, then $\prod_{\pi}(\alpha)$ is a \prod -quantifier-free object for every projection π from I .

Skolem fibration

Theorem (Skolemization)

Every Skolem fibration p validates the principle:

$$\forall x \exists y \alpha(i, x, y) \cong \exists f \forall x \alpha(i, x, fx).$$

Gödel fibration

Definition

A Skolem fibration $p: E \longrightarrow B$ is called a **Gödel** fibration if the sub-fibration $\bar{p}: \bar{E} \longrightarrow B$, whose elements are \prod -quantifier-free objects, has enough \prod -quantifier-free objects.

Gödel fibration

Theorem (Prenex normal form)

In a Gödel fibration $p: E \rightarrow B$, for every element α of a fibre E_I there exists an element β such that

$$\alpha(i) \cong \exists x \forall y \beta(x, y, i)$$

and β is \prod -quantifier-free in the sub-fibration \bar{p} of \prod -quantifier-free objects of p .

The Dialectica fibration

Dialectica construction. Let $p: E \rightarrow B$ be a fibration, whose base category is cartesian closed. Define a category $\mathfrak{Dial}(p)$ as follows:

- **objects** are quadruples (I, X, U, α) where I, X and U are objects of the base category B and $\alpha \in E_{I \times X \times U}$ is an objects of the fibre of p over $I \times X \times U$;
- a **morphism** from (I, X, U, α) to (J, Y, V, β) is a quadruple (f, f_0, f_1, ϕ) where
 - 1 $I \xrightarrow{f} J$ is a morphism in B ;
 - 2 $I \times X \xrightarrow{f_0} Y$ is a morphism in B ;
 - 3 $I \times X \times V \xrightarrow{f_1} U$ is a morphism in B ;
 - 4 $\alpha(i, x, f_1(i, x, v)) \xrightarrow{\phi} \beta(f(i), f_0(i, x), v)$ is an arrow in the fibre over $I \times X \times V$.

Skolemization

Theorem

When the base category \mathcal{B} of a fibration p is cartesian closed, the fibration $\mathcal{D}ial(p)$ satisfies the principle

$$\forall x \exists y \alpha(i, x, y) \cong \exists f \forall x \alpha(i, x, fx)$$

for every α .

Hofstra (2011), *The dialectica monad and its cousins*, Models, logics, and higherdimensional categories: A tribute to the work of Mihály Makkai, 53:107-139

Our contribution

Theorem

Let $p: E \longrightarrow B$ be a fibration with simple products, coproducts and such that B is cartesian closed. Then there exists a fibration p' such that $\mathcal{D}ial(p') \cong p$ if and only if p is a Gödel fibration.

Sketch of the proof

The original Dialectica construction $\mathcal{D}ial$ can be seen as the composition of two free constructions $\mathcal{S}um$ and $\mathcal{P}rod$, which are the simple sum (or co-product) and the simple product completions, respectively.

Lemma

There is an isomorphism of fibrations, natural in p :

$$\mathcal{D}ial(p) \cong \mathcal{S}um(\mathcal{P}rod(p)).$$

These completions are fully dual, in particular $\mathcal{P}rod(p) \cong \mathcal{S}um(p^{op})^{op}$, so we only need to study one and can then deduce results for the other construction.

Hofstra (2011), *The dialectica monad and its cousins*, Models, logics, and higherdimensional categories: A tribute to the work of Mihály Makkai, 53:107-139

Sketch of the proof

Theorem

A fibration $p: E \rightarrow B$ with simple coproducts is an instance of simple coproduct completion if and only if it has enough \coprod -quantifier-free objects. Moreover, in this case $p \cong \text{Sum}(p')$ where p' is the subfibration of \coprod -free-quantifiers objects of p .

Theorem

A fibration $p: E \rightarrow B$ with simple products is an instance of simple product completion if and only if it has enough \prod -quantifier-free objects. Moreover, in this case $p \cong \text{Prod}(p'')$ where p'' is the subfibration of \prod -free-quantifiers objects of p .